US-PAT-NO:

5586264

DOCUMENT-IDENTIFIER:

US 5586264 A

TITLE:

Video optimized media streamer with cache

management

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INVENTOR-INFORMATION:

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US-CL-CURRENT:

725/115, 709/231 , 710/14 , 710/52 , 711/118 ,

725/119

, 725/88 , 725/92 , 725/94 , 725/97

ABSTRACT:

A data storage system includes a mass storage unit storing a data entity,

such as a digital representation of a video presentation, that is partitioned

into a plurality N of temporally-ordered segments. A data buffer is bidirectionally coupled to the mass storage unit for storing up to M of the

temporally-ordered segments, wherein M is less than N. The data buffer has an $\,$

output for outputting stored ones of the temporally-ordered segments. The data $\begin{tabular}{ll} \hline \end{tabular}$

storage system further includes a data buffer manager for scheduling transfers

of individual ones of the temporally-ordered segments between the mass storage $% \left(1\right) =\left(1\right) \left(1\right) \left($

unit and the data buffer. The data buffer manager schedules the transfers in

accordance with at least a predicted time that an individual one of the temporally-ordered segments will be required to be output from the data buffer.

When employed with a media streamer (10) distributed data buffer management

techniques are employed for selecting blocks to be retained in a buffer

memory,

either in a storage node (16, 17) or in a communication node (14). These

techniques rely on the predictable nature of the video data stream, and thus

are enabled to predict the future requirements for a given one of the data blocks.

12 Claims, 26 Drawing figures

Exemplary Claim Number: 1

Number of Drawing Sheets: 14

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Brief Summary Text - BSTX (11):

The playing of movies and video is today accomplished with rather old

technology. The <u>primary storage</u> media is analog tape, such as VHS recorders/players, and extends up to the very high quality and very expensive

 ${\tt D1}\ {\tt VTR's}\ {\tt used}\ {\tt by}\ {\tt television}\ {\tt studios}\ {\tt and}\ {\tt broadcasters}.$ There are many problems

with this technology. A few such problems include: the manual labor required

to load the tapes, the wear and tear on the mechanical units, tape head, and

the tape itself, and also the expense. One significant limitation that troubles Broadcast Stations is that the VTRs can only perform one function at a

time, sequentially. Each tape unit costs from \$75,000 to \$150,000.

Detailed Description Text - DETX (32):

RAID REDUNDANT ARRAY of INEXPENSIVE DISKS: A storage arrangement that uses

several magnetic or optical disks working in tandem to increase bandwidth

output and to provide redundant backup.

Detailed Description Text - DETX (57):

In general, a storage node includes a mass storage unit (or an interface to

a mass storage unit) and a capability to locally buffer data read from or to be

written to the mass storage unit. The storage node may include sequential

access mass storage in the form of one or more tape drives and/or disk drives,

and may include random access storage, such as one or more disk drives accessed

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in a random access fashion and/or semiconductor memory.

Detailed Description Text - DETX (96):

A second advantage of combining hierarchical tape storage to any video

system is that it provides rapid $\underline{\mathbf{backup}}$ to any movie that is stored on disk, in

the event that a disk becomes inoperative. A typical system will maintain a

"spare" disk such that if one disk unit fails, then movies can be reloaded from

tape. This would typically be combined with a RAID or a RAID-like system.

Detailed Description Text - DETX (400):

The buffer operation is managed by the video adapter's controller 226,

placing the N bytes of data in the next available buffer space starting at

address zero of that buffer. Controller 226 keeps track of the length of data

in each buffer and if that data has been "played" or not. Whenever sufficient

buffer space is free, the card accepts the next $\underline{\textbf{WRITE command}}$ and DMA's the

data into that buffer. If not enough buffer space is free to accept the full

data block (typically a Slow Play or Pause condition), the WRITE is not accepted and a buffer full return code is returned.